

Specification and Testing of Protocols in Distributed Object Systems

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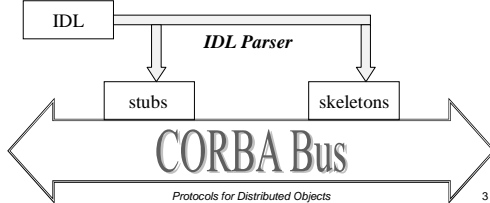
Background: CORBA and IDL

- CORBA
 - Industry standard for distributed object computing
- IDL
 - Notation for defining object interface
 - Programming language neutral

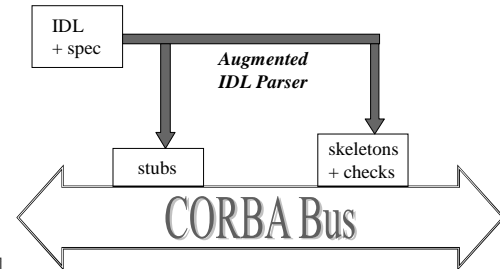


IDL - Interface Definition Language

- IDL description is given to a parser.
- Creates repositories, skeletons, stubs.



Extending the IDL



Example: GUI

```
interface GUI {  
    state bool button_down;  
    state enum {idle, ready, done} status;  
  
    (status == idle) next  
        (status == idle || status == ready)  
    (button_down) next  
        (button_down || status == ready)  
  
    transient.(status == idle && button_down)  
};
```



Challenge

- Temporal logic has high "intimidation factor" for many practitioners
- Focus on a specific subclass of safety properties: protocols
- Design goals:
 - Clean and simple specification notation
 - Automated support for run-time validation



Labeled Transition System

- State
 - Finite enumeration of "phases"
 - Collection of simply typed variables
- Transitions
 - Arcs labeled with messages
 - Method names, argument values
 - Indicate how variables are modified

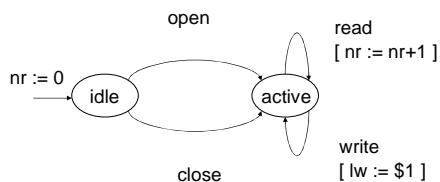


Example: File Access IDL

```
interface FileAccess {
    void open ( );
    void close ( );
    int read ( );
    void write (int);
};
```



Example: File Access



Protocol Specification in CertIDL

```
protocol FileAccess {
    state >idle, active;
    var int nr = 0, lw;
    <idle> open <active>
    <active> read <nr := nr + 1; active>
    <active> write <lw := $1; active>
    <active> close <idle>
};
```



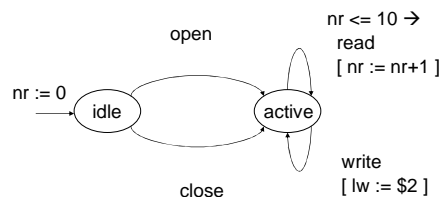
Guarded Transitions

- A transition is enable according to state (phase and variable values)
- "with" clause indicates guard on variables

```
<active> read with nr <= 10
    <nr := nr + 1; active>
```



Example 2: File Access with Guards



Multiplicity of Protocols

- One component may support several protocol *types*

```
protocol FileAccess {...};
protocol UserPrivileges {...};
```
- Each protocol type, may have several *instances*
 - Each instance has its own phase and its own copy of the variables
 - A transition updates the state of a single instance

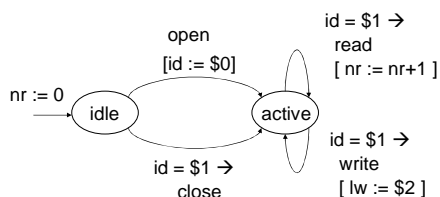


IDL Declaration

```
interface MultipleFileAccess {
    int open ( );
    void close (int);
    int read (int);
    void write (int, int);
};
```



Example 3: Multiple File Access



Protocol Specification in CertIDL

```
protocol MultipleFileAccess {
    state >idle, active;
    var int nr = 0, lw, id;
    <idle> open <id := $0; active>
    <active> read with $1 = id
        <nr := nr + 1; active>
    <active> write with $1 = id
        <lw := $2; active>
    <active> close with $1 = id <idle>
};
```



Protocol Types: Shared State and Transitions

- Some state is common to all instances
 - eg, the number of open files
 - "shared state": phases and variables
- Instance transitions:
 - Can use shared state in guards
 - Can not modify shared state
- Shared transitions
 - Can not use or modify instance state
 - Enabled when an instance can accept



Protocol Specification in CertIDL

```
protocol MultipleFileAccess {
    state local >idle, active;
    var local int nr = 0, lw, id;
    var shared int n = 0;

    open <n := n+1>
    close <n := n-1>
    <idle> open <id := $0; active>
    <active> read with $1 = id && n <= 10
        <nr := nr + 1; active>
    <active> write with $1 = id
        <lw := $2; active>
    <active> close with $1 = id <idle>
};
```



Singleton Protocols

- Some protocols are not meant to be multiply instantiated
 - Original FileAccess protocol
 - As written, would accept:
 - open read open close close
- "Singleton"
 - No shared state or shared transitions



Singleton Protocol

```

singleton protocol FileAccess {
  state local >idle, active;
  var local int nr = 0, lw;
  <idle> open <active>
  <active> read <nr := nr + 1; active>
  <active> write <lw := $1; active>
  <active> close <idle>
};
    
```

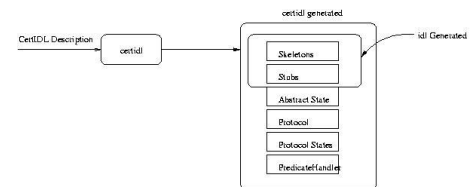


Other Features

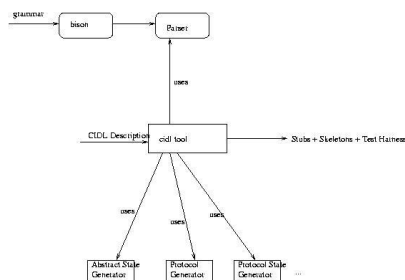
- Synchronized transitions
 - Multiple instances accept single message
- Global state
 - Abstract state, implementation provided
- Synchronous / asynchronous support
- Messages received and sent



CertIDL

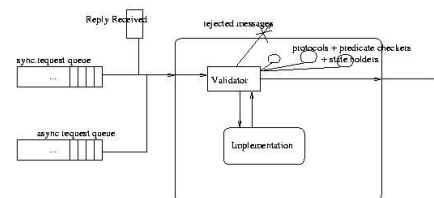


Tool Architecture

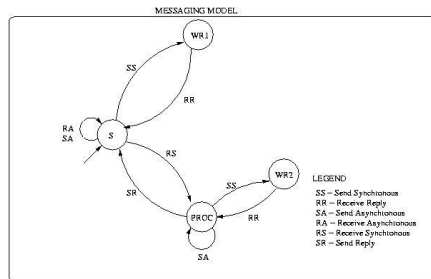


Run-Time Validation

RUN TIME MONITORING



Synchrony vs Asynchrony



Related Work

- OO community
- CFSM (S, Prospec)
- Petri nets (P-nut, PROTEAN)
- Estelle (NBS)
- LOTOS (Sedos)
- SDL



Distinguishing Features

- Design focus: run-time validation
 - Testing harness implementation
- Hierarchy of protocol "granularity"
 - Instance, class, abstract state
 - Multiplicity of instantiation
 - Coupling transitions, sharing state
- Inclusion of message arguments



Current Status

- Specification notation
 - Limitations in mixing send/receive, and synchronous/asynchronous messages
- Implementation
 - CertIDL language and tool suite extended to support protocol validation
- Evaluation
 - Industrial scale application: phone number activation



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