

Deadlock, Livelock, and Starvation



Outline

- **Situations preventing packet delivery**
- **Theory of Deadlock Avoidance**
 - Channel Dependency
 - Necessary and sufficient condition
 - Relaxing this constraint
 - Extended Channel Dependency Graph
- **Deadlock in MINs**

Situations preventing packet delivery

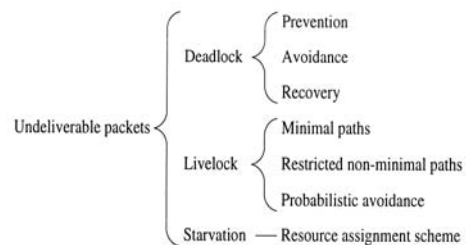


Figure 3.1: A classification of the situations that may prevent packet delivery

Theory of Deadlock Avoidance

- **Basic definitions**
 - **routing function** (set of output channels based on the current and destination nodes)
 - **selection function** (one free channel is selected out of many, if available)
 - **configuration** (assignment of a set of packets or flits to each queue) based on a *reachable* or *routable* assignment
 - legal
 - illegal
 - **connected routing function** (able to establish a path between each pair of nodes)

Deadlock Configuration

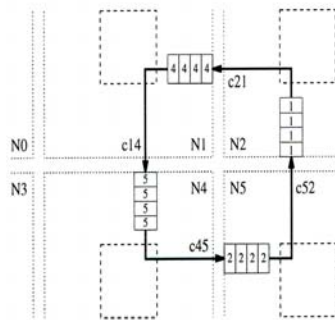


Figure 3.2: Deadlocked configuration for R

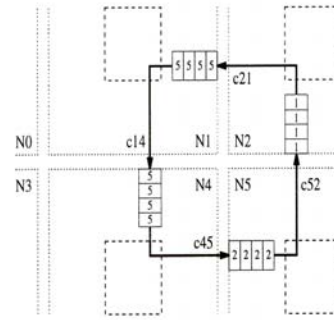


Figure 3.3: Illegal configuration for R

Avoiding Deadlock

- Identification of *channel dependency* under a routing function R
- Necessary and Sufficient Condition
 - Breaking the dependency by providing some restriction on R

Breaking Channel Dependency with Virtual Channels on a Ring

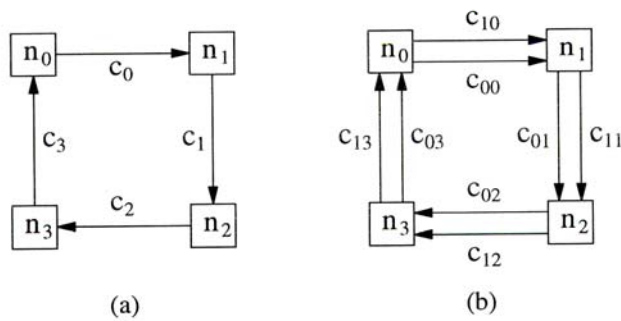


Figure 3.4: Networks for the example

Can we Avoid Deadlock in the Presence of Cyclic Dependency?

- No until 1993
- All networks were built with the notion of breaking cyclic dependency
- Prof. Duato provided a new framework
 - concept of escape channel
 - routing subfunction corresponding to escape channels should not have cyclic dependency
 - has led to flexible routing schemes
 - led to simpler adaptive routing schemes

Deadlock Avoidance in the Presence of Cyclic Dependency

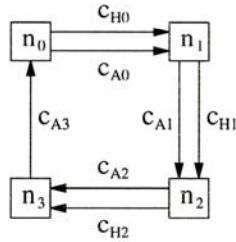


Figure 3.5: Network for the example

Use either C_{A_i} ($j \neq i$) or C_{H_i} $j > i$

Extended Channel Dependency Graph

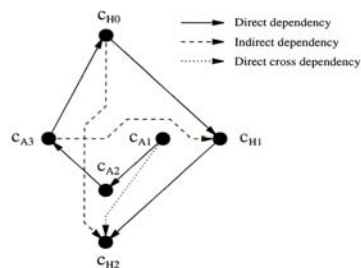


Figure 3.8: Extended channel dependency graph for R_1 using wormhole switching

R1: use C_{A_i} if $j < i$ or C_{H_i} if $j > i$

Developing Adaptive Routing Algorithms

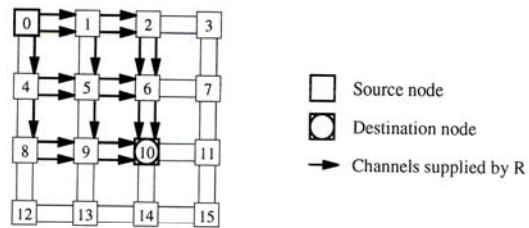


Figure 3.9: Routing example for R

Deadlock in MINs

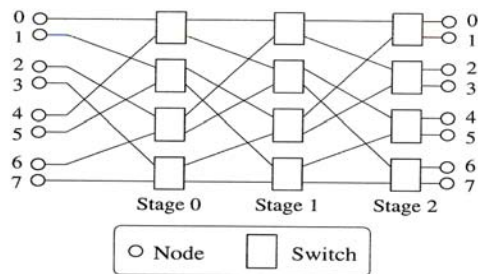


Figure 3.13: A multistage interconnection network

Deadlock Recovery

- Recently has generated some interest
- Deadlock Avoidance restricts some use of channels/buffers
- Can we trade-off between better resource utilization and deadlock recovery?
- Reserve minimal resources for recovery
- Should limit which packets in deadlock are eligible for recovery

Deadlock Recovery (cont.)

- Recovery Process
 - Progressive
 - Regressive
- Deadlock detection vs. blocking
- Heuristics for deadlock detection
 - time-outs (value dependent on message length)
 - may not detect deadlock for a long time or
 - may “detect” deadlocks that do not exist
- not applied to any real system yet

Livelock

- **Occurs if a packet circulates forever in the network**
- **Limiting misrouting avoids livelock**
- **Minimal routing algorithms allow no misrouting**
- **Misrouting is useful for fault tolerance**
- **Deflection routing (any free channel) provides probabilistic livelock freedom**